

## ON SHORT EDGES IN STRAIGHT-EDGE TRIANGULATIONS

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### Abstract.

If a triangulation is drawn with straight edges, the ratio of the lengths of the shortest and the longest edges does not have to go to zero, even if the number of vertices goes to infinity. In this paper bounds are given for the above ratio, when certain restrictions are placed on the maximum degree of the triangulation.

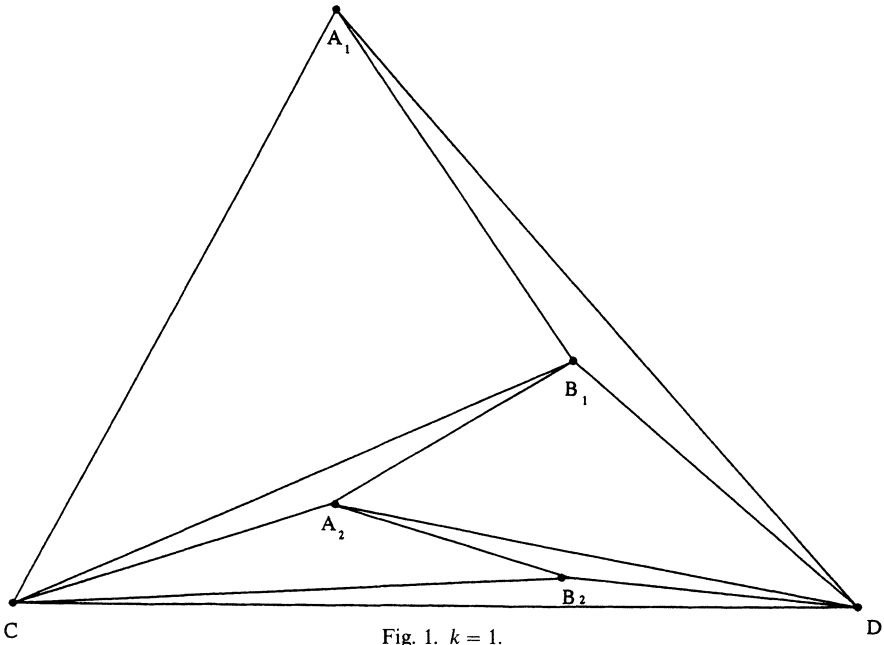
A triangulation of the plane is a planar graph in which each region, including the infinite one, is bounded by exactly three edges. For standard terminology in graph theory, see, for example, [1]. The triangulation is said to be straight-edged if the graph is drawn with line segments as edges, so that it represents a triangle which is subdivided into smaller triangles. In this paper, a triangulation is always taken to mean a straight-edged triangulation.

Denote by  $T_n$  a triangulation with  $n$ -vertices and let  $d(T_n)$  and  $l(T_n)$  respectively by the lengths of the longest and shortest edges of  $T_n$ . The ratio  $l(T_n)/d(T_n)$  is denoted by  $\alpha(T_n)$ .

As the number of vertices in a triangulation increases, the area of the smallest triangle divided by the area of the outer triangle must decrease to zero. However, such is not the case with the length of the shortest edge.

We constructed triangulations  $T_n$ , with  $n$  arbitrarily large satisfying  $\alpha(T_n) > 1/4$ . Two persons at a seminar of Micha A. Perles in Jerusalem ([2]), after seeing our construction, provided a better one. Unfortunately, their names are not known.

Figure 1 is an example of such a triangulation. The large triangle has vertices  $C$  at  $(0, 0)$ ,  $D$  at  $(3, 0)$  and  $A_1$  on  $x = 1$ .  $B_1$  is any point on  $x = 2$  within  $A_1CD$  and it is joined to  $A_1$ ,  $C$  and  $D$ .  $A_2$  is any point on  $x = 1$  within  $B_1CD$ , and it is joined to  $B_1$ ,  $C$  and  $D$ . The point  $B_2$  is chosen in a similar manner. The construction can be continued indefinitely. The length  $l(T_n)$  is greater than 1 regardless of how large

Fig. 1.  $k = 1$ .

$n$  becomes. By compressing the large triangle towards  $CD$  if necessary,  $CD$  will be the longest edge. Hence,  $d(T_n) = 3$  and  $\alpha(T_n) > 1/3$ .

Note that the above bound is the sharpest possible. Here is a sketch of the proof:

Assume that  $d(T_n) = 1$ , so that the area of the large triangle is at most  $1/2$ . It is divided into exactly  $2n-5$  smaller triangles. It follows that at most  $c\sqrt{n}$  of them can have an area exceeding  $1/\sqrt{n}$ , where  $c$  is a positive constant. All others, that is at least  $2n-5-c\sqrt{n}$ , have one angle very close to  $\pi$  and two angles very close to  $0$  when  $n$  is large and  $\alpha(T_n)$  remains above a constant value (since the area of a triangle with sides  $a$  and  $b$  and angle  $\theta$  between them is  $(ab \sin \theta)/2$ ).

Among these triangles, one can find two with a common edge, say  $xvy$  and  $yvz$ , such that angle  $xvy$  is very close to  $\pi$  and angle  $yvz$  is very close to  $0$ . Either  $x-v-y-z$  or  $x-v-z-y$  is almost a straight line. Since  $d(T_n) = 1$ , we cannot have  $l(T_n) > 1/3 + \varepsilon$  for any fixed positive  $\varepsilon$  and large  $n$ .

In Figure 1, both  $C$  and  $D$  are joined to all other vertices as well as to each other. It is natural to ask what bounds can be obtained for  $\alpha(T_n)$ , when certain restrictions are placed on  $\Delta(T_n)$ , the maximum degree of  $T_n$ .

**THEOREM A.** *Let  $k$  be a fixed positive integer. Then, for an infinite number of values of the integer  $n$ , there exists a triangulation  $T_n$  which satisfies*

$$\alpha(T_n) > 1/(2k + 1), \text{ with } \Delta(T_n) < 2(n/2)^{1/k} + 2.$$

PROOF. Let  $n = 1 + 2(1 + (m - 1) + \dots + (m - 1)^k)$ , where  $m$  is a positive integer; note that  $2(m - 1)^k < n$ . A triangulation  $T_n$  with the desired properties is constructed.

The argument is inductive. The case  $k = 1$  has already been dealt with in Figure 1. To give a clearer illustration of the general construction, consider the case  $k = 2$ . Figure 2 illustrates the case  $m = 2$ .

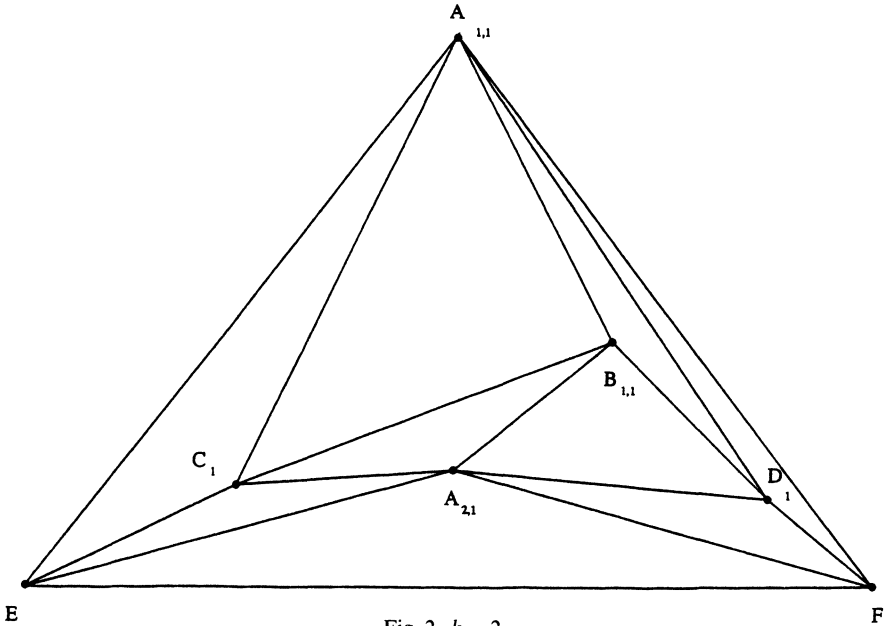


Fig. 2.  $k = 2$ .

Here, the large triangle has vertices  $E$  at  $(0, 0)$ ,  $F$  at  $(5, 0)$  and  $A_{1,1}$  on  $x = 2$ . Take  $C_1$  on  $x = 1$  and  $D_1$  on  $x = 4$  such that  $C_1D_1$  is parallel to  $EF$  and lies within  $A_{1,1}EF$ . Join  $C_1$  to  $E$  and  $A_{1,1}$ , and join  $D_1$  to  $F$  and  $A_{1,1}$ . We now triangulate  $A_{1,1}C_1D_1$  as in the case  $k = 1$ , using the points  $A_{1,2}A_{1,3}, \dots, A_{1,m-1}$  on  $x = 2$  and the points  $B_{1,1}, B_{1,2}, \dots, B_{1,m-1}$  on  $x = 3$ . Note, however, that  $C_1$  is not joined to  $D_1$ .

Take  $A_{2,1}$  on  $x = 2$  within  $B_{1,m-1}C_1D_1$  and join it to all three vertices as well as  $E$  and  $F$ . Take  $C_2$  on  $x = 1$  and  $D_2$  on the line  $x = 4$  such that  $C_2D_2$  is parallel to  $EF$  and lies within  $A_{2,1}EF$ . The construction is repeated until  $A_{m,1}$  is chosen and joined to  $E$  and  $F$ .

The total number of vertices is exactly equal to  $n$ . The degree of  $A_{1,1}$  (and of  $A_{m,1}$ ) is 5. The degree of  $A_{i,1}$  is 8 for  $2 \leq i \leq m - 1$ . The degree of  $A_{i,j}$  (and of  $B_{i,j}$ )

is 4 for all other values of  $i$  and  $j$ . The degree of  $C_i$  (and of  $D_i$ ) is  $2m$  for all  $i$ , as is the degree of  $E$  (and of  $F$ ). Hence,  $\Delta(T_n) = 2m < 2\sqrt{n/2} + 2$ . By compressing the large triangle towards  $EF$  if necessary, we can make  $EF$  the longest edge. Hence,  $d(T_n) = 5$ . Since  $l(T_n) > 1$ ,  $\alpha(T_n) > 1/5$ .

We now consider the general case. Let the large triangle be  $A_{1,1}EF$ , with  $E$  at  $(0, 0)$ ,  $F$  at  $(2k + 1, 0)$  and  $A_{1,1}$  on  $x = k$ . Take  $C_1$  on  $x = 1$  and  $D_1$  on  $x = 2k$  such that  $C_1D_1$  is parallel to  $EF$  and lies within  $A_{1,1}EF$ . Join  $C_1$  to  $E$  and  $A_{1,1}$ , and join  $D_1$  to  $F$  and  $A_{1,1}$ . We can triangulate  $A_{1,1}C_1D_1$  as in the case  $k - 1$ . We now take an appropriate point  $A_{2,1}$  on  $x = k$  and continue with the inductive construction.

The total number of vertices in this triangulation is exactly  $n$ . It is routine to verify that, as in the case  $k = 2$ , none of the vertices has a degree exceeding  $2m < 2(n/2)^{1/k} + 2$  and that  $\alpha(T_n) > 1/(2k + 1)$ . This completes the proof of Theorem A.

**THEOREM B.** *For fixed positive integer  $k$  and  $n$  sufficiently large, every  $T_n$  satisfies  $\alpha(T_n) < 1/k$ , provided that  $\Delta = \Delta(T_n) < (\pi n/6k^2)^{1/(k+1)}$ .*

**PROOF.** The proof uses an indirect argument. Consider any  $T_n$  with  $\Delta(T_n)$  as provided. Assume that  $d(T_n) = 1$ , so that  $l(T_n) = \alpha(T_n)$ . Suppose that  $\alpha(T_n) \geq 1/k$ , and derive a contradiction.

We call a triangle in  $T_n$  good if its area is less than  $S/2$ , where  $S = 5\Delta^k/n$ . Note that  $S$  is very close to 0 when  $n$  is sufficiently large.

Call an angle  $\theta$  thin if it satisfies  $0 < \theta < 6/5 k^2 S$  and call it thick if  $\pi - 6/5 k^2 S < \theta < \pi$ .

Let  $a$  and  $b$  be the lengths of two edges of a good triangle, and let  $\theta$  be the angle between them. Then  $ab \sin \theta < S$ . Since we assume that  $a \geq 1/k$  and  $b \geq 1/k$ ,  $\sin \theta < k^2 S$ . Since  $k$  is fixed,  $\sin \theta$  is very close to 0 and it follows that  $\theta$  is either thin or thick. Thus each good triangle has exactly one thick and two thin angles.

Consider a vertex  $v$  surrounded by good triangles. Then every angle formed by two consecutive edges at  $v$  in clockwise order is either thick or thin, with at most two thick angles formed by consecutive edges at  $v$ . The sum of all the thin angles formed by consecutive edge at  $v$  is less than  $\Delta(6/5)k^2 S < \pi$ . Hence, there are exactly two thick angles formed by consecutive edges at  $v$ . It follows that for any edge incident with  $v$ , there is another edge incident with  $v$  and forming a thick angle with the given edge.

Note that the two edges forming a thick angle are not required to lie on the same triangle of  $T_n$ .

Let  $V$  denote the set of vertices not surrounded by good triangles. In particular, the three vertices of the large triangle are in  $V$ . Since the area of this triangle is at most  $1/2$ , there are at most  $1/S$  triangles which are not good, so that  $|V| \leq 3(1 + 1/S)$ . The total number of vertices accessible from  $V$  by a path of

length not exceeding  $k$  edges is at most  $3(1 + 1/S)(1 + \Delta^2 + \dots + \Delta^k) < 5\Delta^k/S = n$ , since  $\Delta \geq 3$  and  $S$  is small.

This means that there exists a vertex  $x_0$  which is inaccessible from  $V$  by any path of length not exceeding  $k$  edges. In particular,  $x_0$  is surrounded by good triangles. Let  $x_1$  be any vertex adjacent to  $x_0$ . Then  $x_1$  is also surrounded by good triangles. Hence, among vertices adjacent to  $x_1$ , we can choose  $x_2$  such that angle  $x_0x_1x_2$  is thick. We can continue to choose vertices in this manner until we have  $x_{k+1}$ , which may no longer be surrounded by good triangles.

By our assumption  $x_i x_{i+1}$  has length at least  $1/k$ , for  $0 \leq i \leq k$ . Thus, the distance between  $x_0$  and  $x_{k+1}$  is at least  $(k + 1 - \frac{1}{2})/k$ . It follows that the distance between  $x_0$  and  $x_{k+1}$  exceeds 1, contradicting  $d(T_n) = 1$ . This completes the proof of Theorem B.

Our results are open for improvements. For example, Theorem A shows by example that  $\alpha(T_n) > \frac{1}{3}$  and  $\Delta(T_n) < 1.42n^{1/2}$  is possible, whereas Theorem B shows that  $\alpha(T_n) > \frac{1}{3}$  and  $\Delta(T_n) < 0.52n^{1/6}$  is impossible. Thus there are gaps in the size of  $\Delta(T_n)$  for which the situation is still unresolved.

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#### REFERENCES

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2. M. Perles, *Private communication*.

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